

Lecture 10: LR parsing and resolving conflicts

- **What are conflicts**
- **Example 1: a simple, unambiguous grammar**
 - **ocamlyacc output**
 - **Using parse trees to understand conflict**
 - **Fixing conflict**
- **Example 2: ambiguous grammar for conditional expressions**
 - **Eliminating conflicts using %prec declarations**

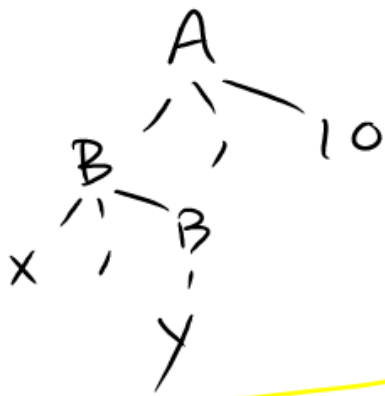
Conflicts

- **ocamlyacc generates tables saying what action to take at each point in parse**
 - **Method is called “LALR(1)”**
 - **“LR(1)” is a similar, but somewhat more powerful, method — will often use “LR(1)” and “LALR(1)” as synonyms.**
- **Not every grammar can be parsed using this method.**
 - **Problem is *always* that ocamlyacc cannot decide on the proper action in some cases**
 - **“Shift/reduce conflict” — cannot decide whether to shift or reduce**
 - **“Reduce/reduce conflict” — knows to reduce, but can’t decide which production to use**

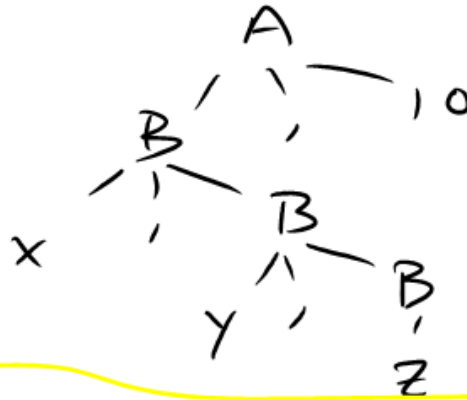
Example 1

- $A \rightarrow B, int$
 $B \rightarrow id \mid id, B$
- **Grammar is unambiguous, but consider these two inputs:**
 - **x,y,10**
 - **x,y,z,10**
- **Both lead to an identical stack/lookahead configuration, but the correct action in one case is shift and in the other is reduce.**
- **Look at s-r parse, and at two parse trees.**

x, y, 10



x, y, z, 10



A → B, int
 B → id
 id, B

Sh		x, y, 10
Sh	x	, y, 10
Sh	x,	y, 10
R	x, y	, 10
R	x, B	, 10
	B	, 10

Sh		x, y, z, 10
Sh	x	, y, z, 10
Sh	x,	y, z, 10
Sh	x, y	, z, 10
R	x, y, z	, 10
	x, y, B	, 10

Example 1 (cont.)

- Presented to `ocamlyacc`:

```
%token int id comma
%start A
%type <int> A
%%
A: B comma int {0}
B: id {0}
  | id comma B {0}
```

- Using "`ocamlyacc -v`", file `simple.output` contains:

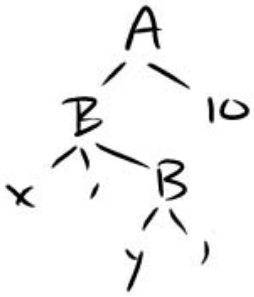
```
3: shift/reduce conflict (shift 6, reduce 2) on comma
state 3
B : id . (2)
B : id . comma B (3)

comma shift 6
```

Example 1 (cont.)

- One way to fix grammar:

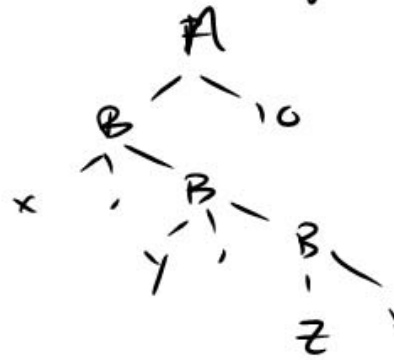
$A \rightarrow B \mid \text{int}$
 $B \rightarrow \text{id}, \mid \text{id}, B$



$S_h \times 2$
 S_h

$x,$
 $x, y,$

$x, y,$



x, y, z, id
 y, z, id
 $, z, \text{id}$

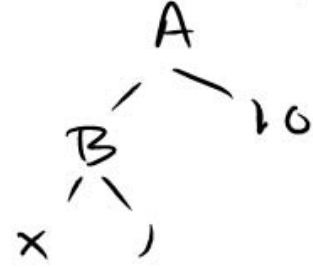
x, y, z, id

if id on stack: shift

if id + , on stack:

if lookahead is id: shift

if number: reduce
if comma: reject



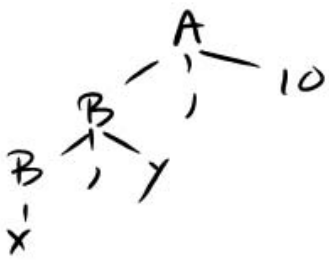
$x, , z, \text{id}$

Example 1 (cont.)

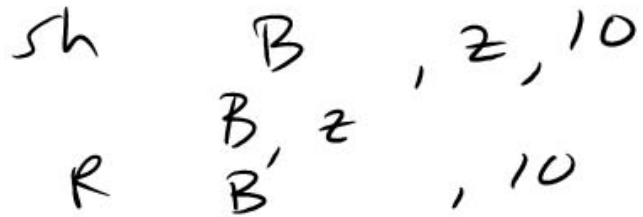
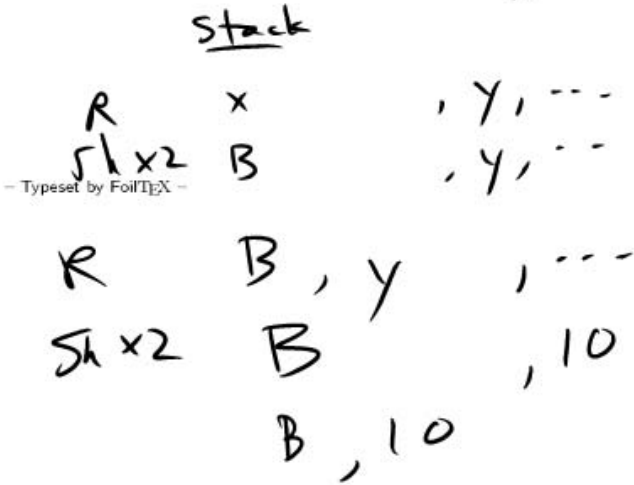
- Another way to fix grammar:

$A \rightarrow B, int$

$b \rightarrow id | B, id$



Stack + lookahead symbol gives enough info. to take correct parse action



Example 2

- **Ambiguous grammar for conditional expressions:**

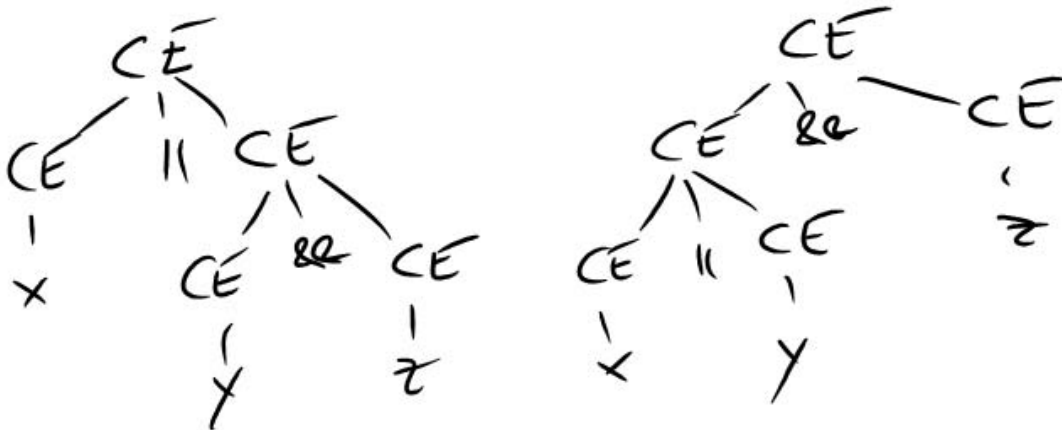
- $CondExpr \rightarrow id \mid CondExpr \parallel CondExpr$
 $\mid CondExpr \&\& CondExpr \mid ! CondExpr$

- **Consider this input:**

- $x \parallel y \&\& z$

- **Leads to a stack/lookahead configuration in which shifting and reducing both work, but produce different parse trees.**

- **Look at s-r parse, and at two parse trees.**



x || y & z

Stack

Input

sh

x || y & z

R

x

|| y & z

sh x 2

CE

|| y & z

R

CE || y

& z

sh x 2

CE || CE

& z

R

CE || CE & z

sh x 2 → R

CE

& z

R

CE || CE & CE

R

CE & z

cof

R

CE || CE
CE

R

CE & CE
CE

Example 2 (cont.)

- **ocamlyacc -v output contains:**

10: shift/reduce conflict (shift 7, reduce 2) on and

10: shift/reduce conflict (shift 8, reduce 2) on or

state 10

~~CondExpr : CondExpr or CondExpr (2)~~

~~CondExpr : CondExpr or CondExpr . (2)~~

~~CondExpr : CondExpr . and CondExpr (3)~~

and shift 7

or shift 8

\$end reduce 2

Example 2 (cont.)

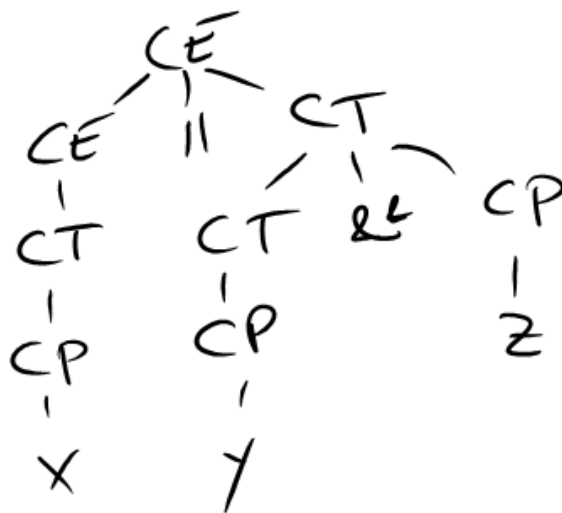
- One way to resolve conflict: fix grammar.
- Use “stratified grammar,” as for arithmetic expressions:

$CondExpr \rightarrow CondTerm \mid CondExpr \parallel CondTerm$

$CondTerm \rightarrow CondPrimary \mid CondTerm \&\& CondPrimary$

$CondPrimary \rightarrow id \mid ! CondPrimary$

$x \parallel y \&\& z$



Example 2 (cont.)

- Another way to resolve conflict: precedence declarations.
- Suppose t_1 is the topmost terminal symbol on the stack, and t_2 is the lookahead symbol. Then:
 - If t_1, t_2 appear in the same %left declaration, then reduce
 - If t_1, t_2 appear in the same %right declaration, then shift
 - If t_1 appears in a declaration before t_2 , shift
 - If t_1 appears in a declaration after t_2 , reduce

%left token, .. CE || CE || ..
%right token, --- CE || CE || ..
%nonassoc token, --- CE || (CE || CE)

Example 2 (cont.)

- Use the ambiguous grammar, but add these declarations:

```
%left or
%left and
```

- $x \ || \ y \ \&\& \ z$ is now handled correctly.

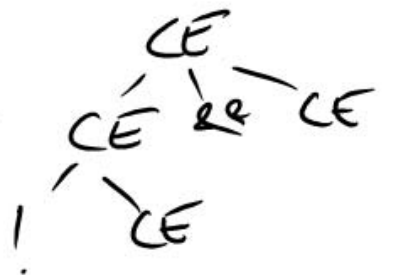
x	$\ \ y \ \&\& \ z$
$C\bar{E}$	$\ \ y \ \&\& \ z$
$C\bar{E} \ \ y$	$\ \&\& \ z$
$C\bar{E} \ \ C\bar{E}$	$\ \&\& \ z$
$C\bar{E} \ \ C\bar{E} \ \&\&$	$\ t$

Example 2 (cont.)

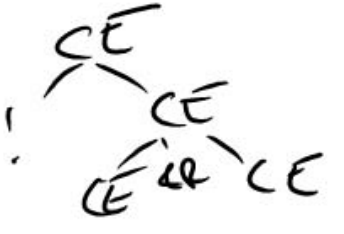
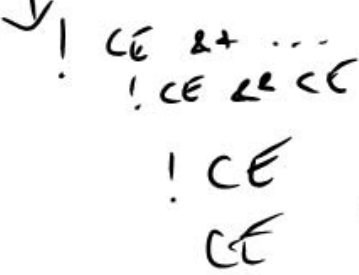
- However, `ocamlyacc` still reports conflicts. Verbose output:

```

6: shift/reduce conflict (shift 7, reduce 4) on and
6: shift/reduce conflict (shift 8, reduce 4) on or
state 6
CondExpr : CondExpr . or CondExpr (2)
CondExpr : CondExpr . and CondExpr (3)
CondExpr : not CondExpr . (4)
and shift 7
or shift 8
$end reduce 4
  
```



- Problem is that we didn't resolve ambiguity involving "!". Add "`%nonassoc not`" after above two lines.



More on conflicts

- **Posted supplementary notes discuss four grammars that have conflicts, and how to resolve them.**
- **All are relevant to the current MP.**